Correctness and Performance for Stateful Chained Network Functions

Junaid Khalid^{W,G} and Aditya Akella^W





Hardware NF → software NF over commodity server

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Intrusion detection system (IDS)

Hardware NF → software NF over commodity server



Intrusion

detection

system (IDS)



Caching proxy

Hardware NF → software NF over commodity server



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Firewall

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WAN optimizer

Hardware NF → software NF over commodity server

- Enables resource consolidation
- **Dynamic allocation** of packet processing
- Adding new functionality



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Hardware NF → software NF over commodity server

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- **Dynamic allocation** of packet processing
- Adding new functionality
- Simplifies service chaining



Intrusion

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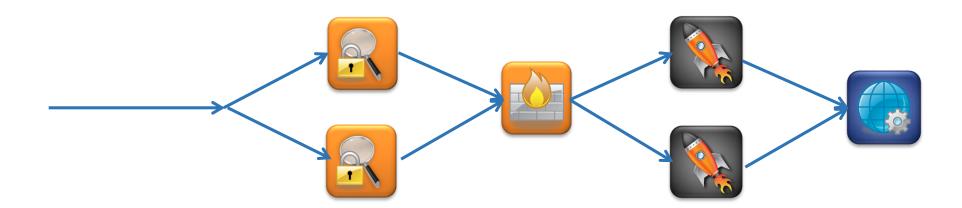
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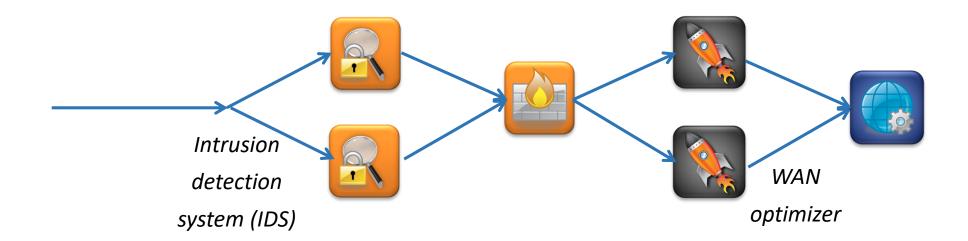


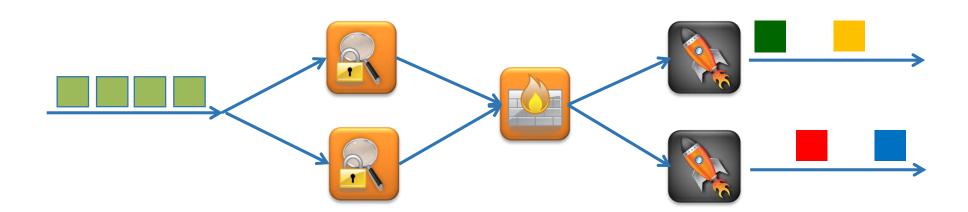


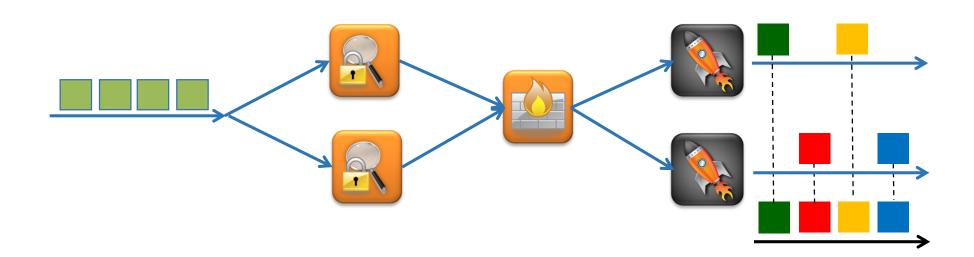


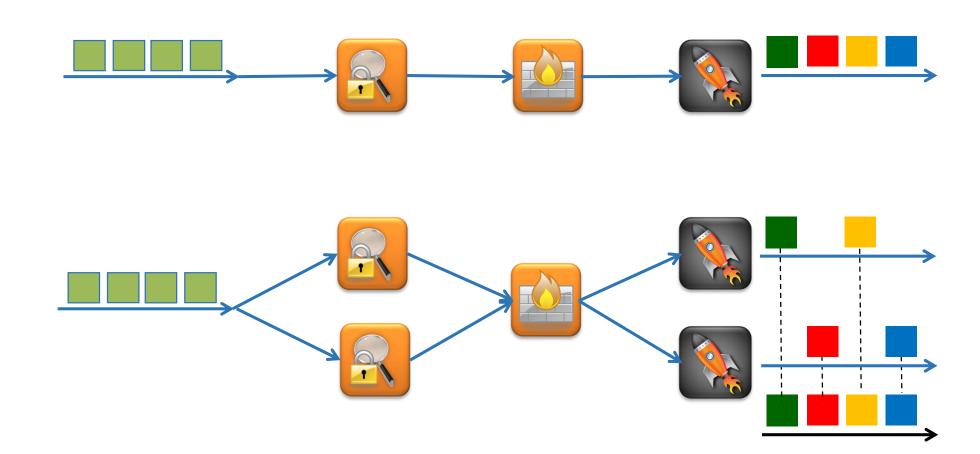




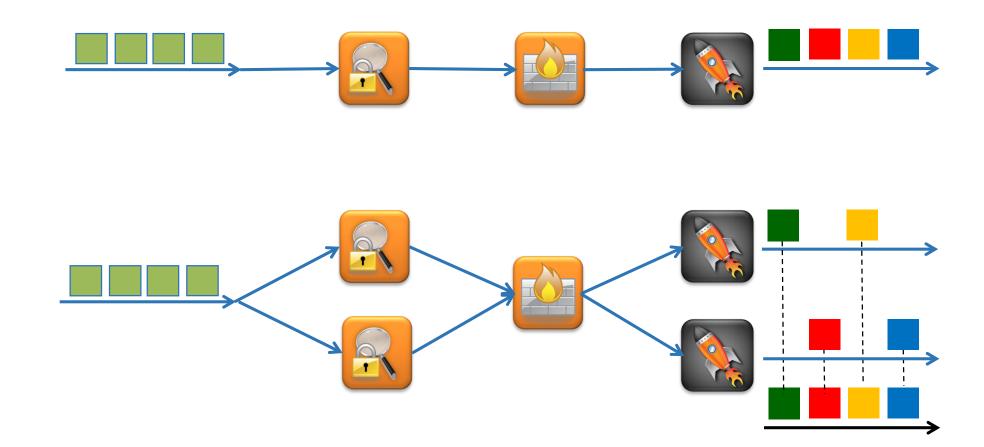








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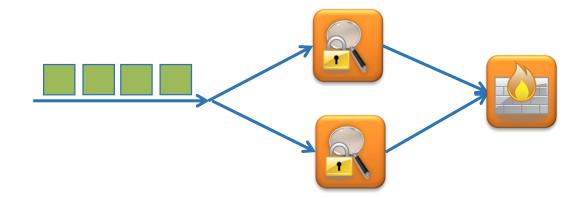
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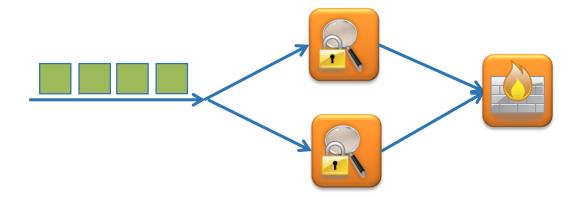
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Ensuring COE is challenging: NF chain attributes & Dynamic Actions



1. NF statefulness

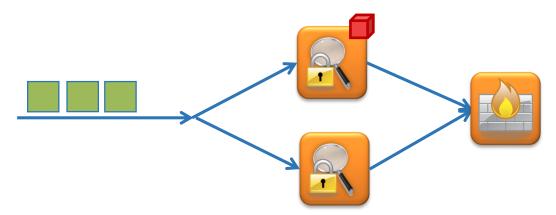
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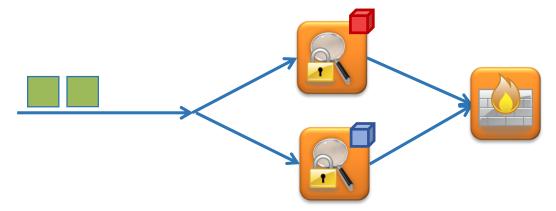
IDS maintains *cross-flows state* (e.g., per host active conn. count) and *per-flow state* (e.g., TCP conn. state)



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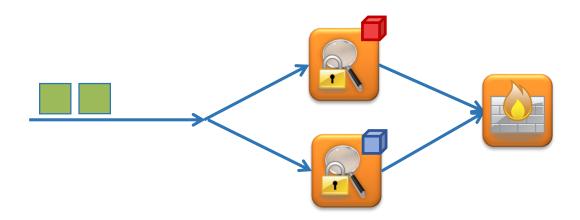
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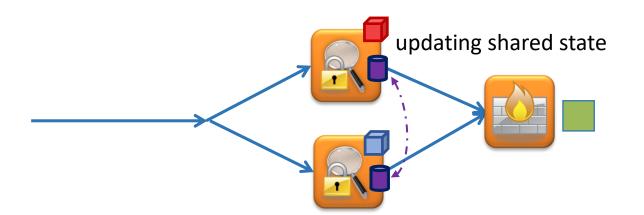
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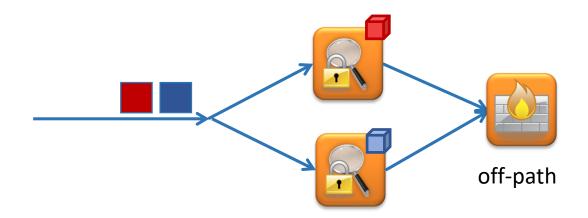
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- 2. Consistent state updates

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- Action taken by an NF instance depends on the state updates from other NF instances



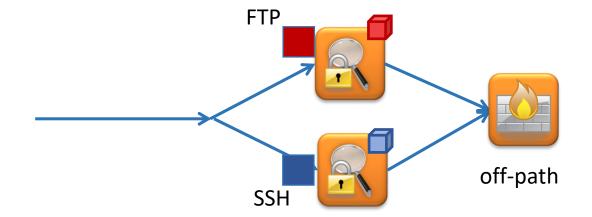
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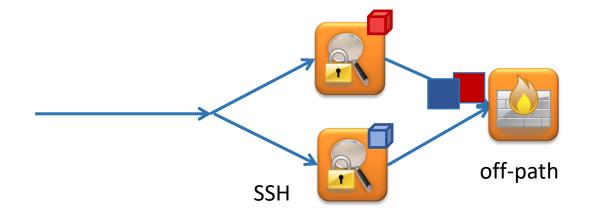
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- 3. Dependency between different NF instances

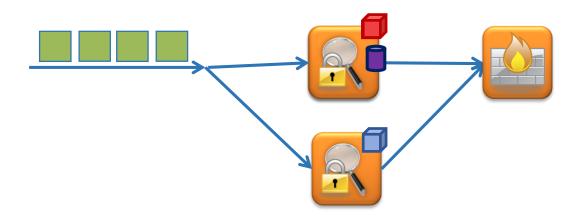
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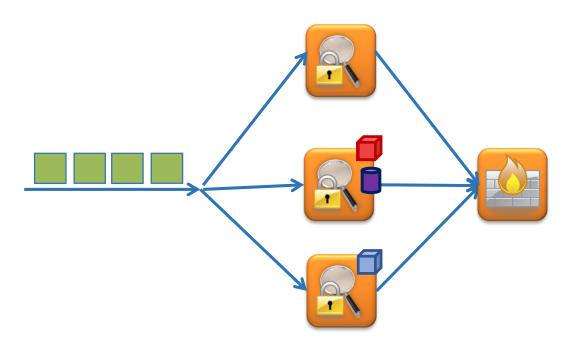
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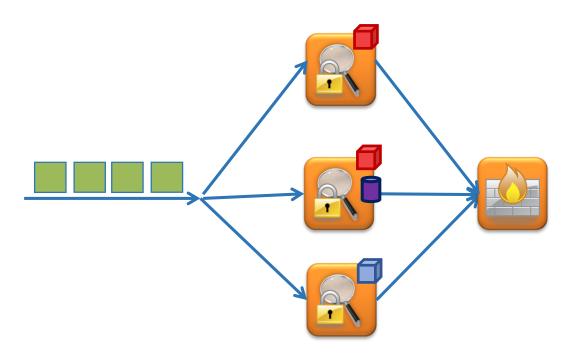
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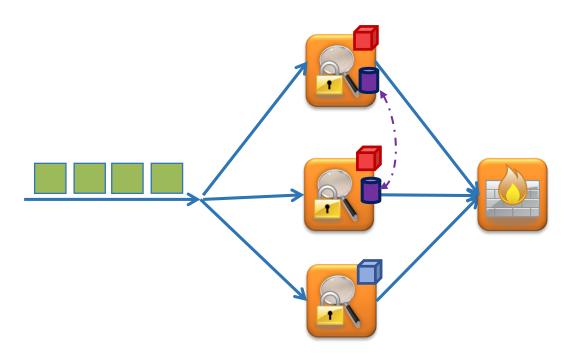


Key requirements

 Safe cross-instance state transfer

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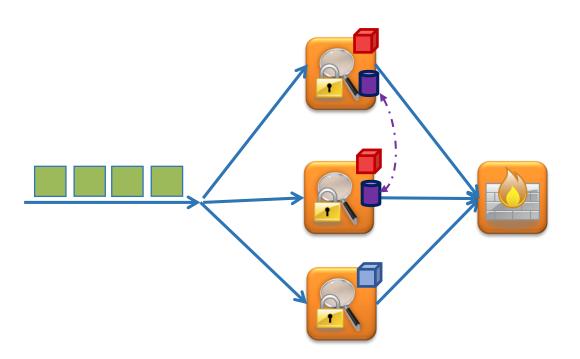


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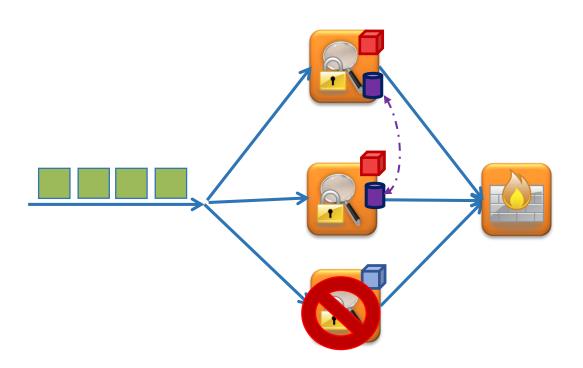
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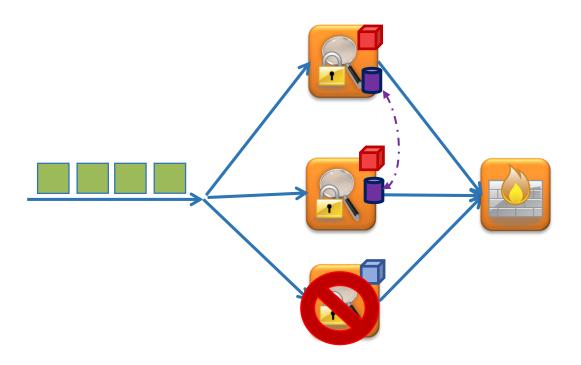
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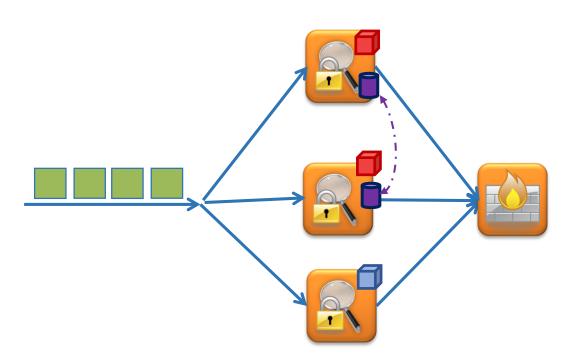
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Failure recovery

• When NF fails, all its state disappears. For fault tolerance, that state needs to be recovered



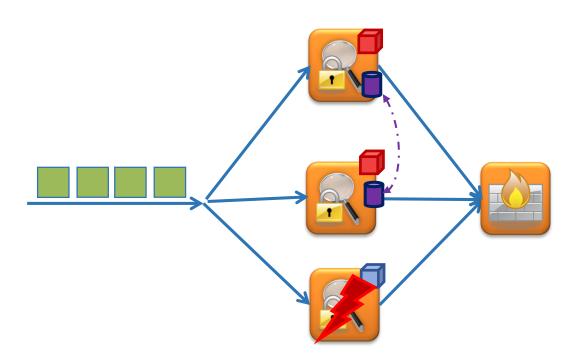
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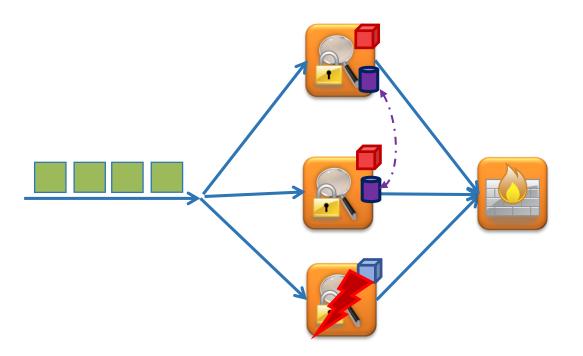
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- Safe cross-instance state transfer
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Instance slowdown

- Clones may be launched to handle a straggler NF (a slow NF)
- Downstream NFs rely on the order at upstream NFs



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Existing Solutions

Framework	State availability	State transfer	Consistent shared state	Duplicate suppression	Chain-wide ordering
Split/Merge[NSDI'13]					
OpenNF[sigcomm'14]					
FTMB [SIGCOMM' 15]					
S6 [NSDI'18]					
Pico Rep.[socc'13]					
StatelessNF[NSDI'17]					

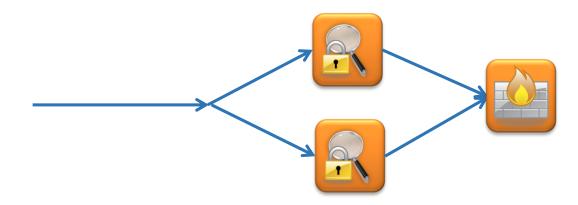
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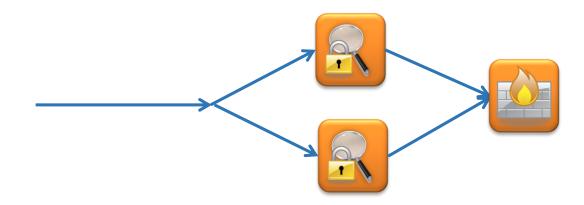
Incomplete support → restricted functionality

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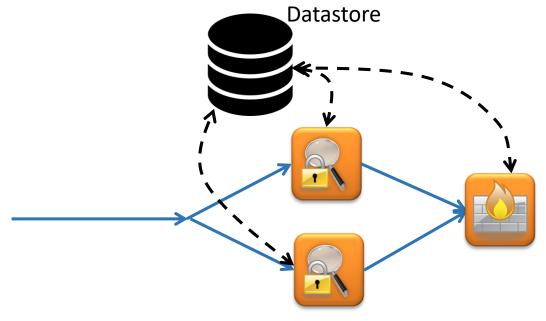
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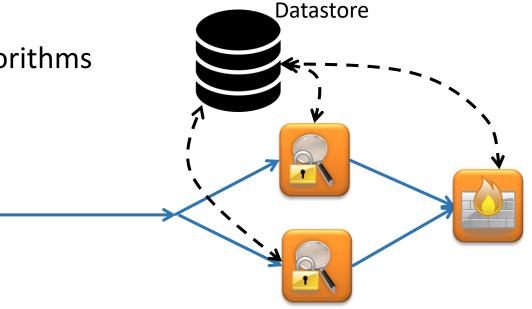
CHC consist of three main building blocks

1. State store external to NF



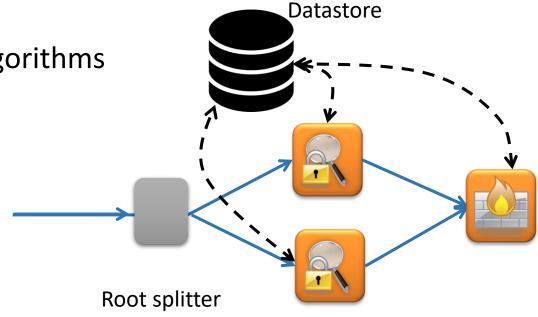
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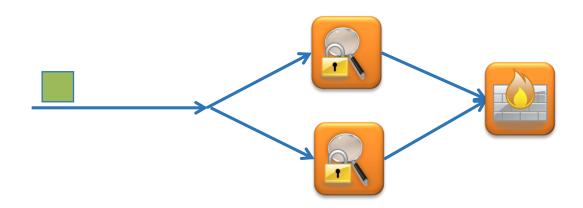
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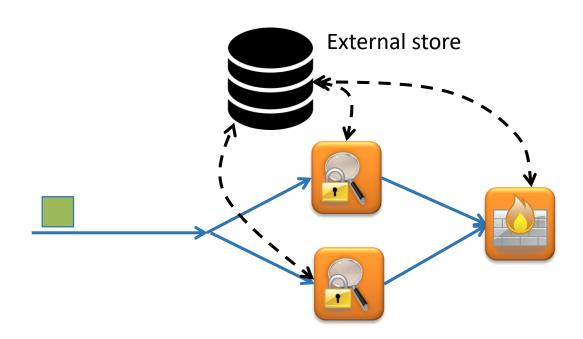
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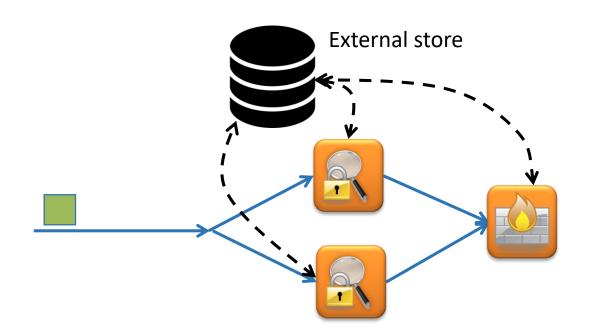
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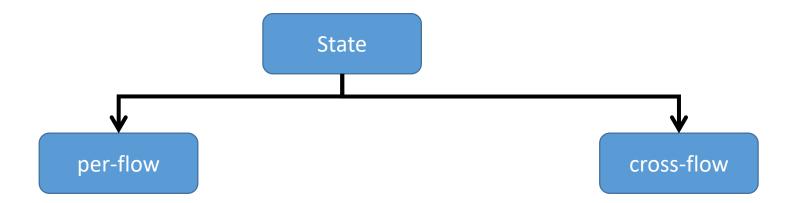


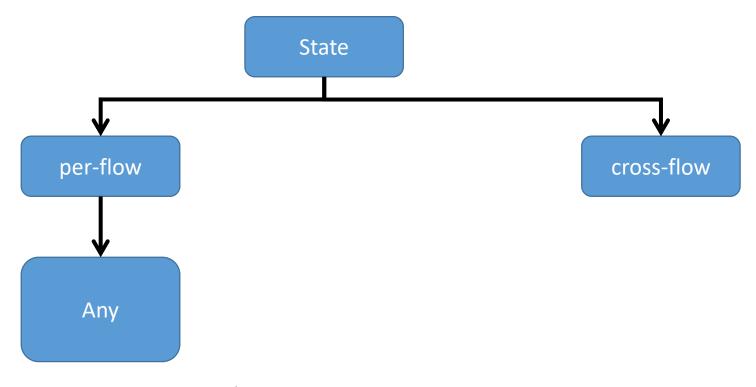
Naively *externalizing* the state can *degrade* NF performance

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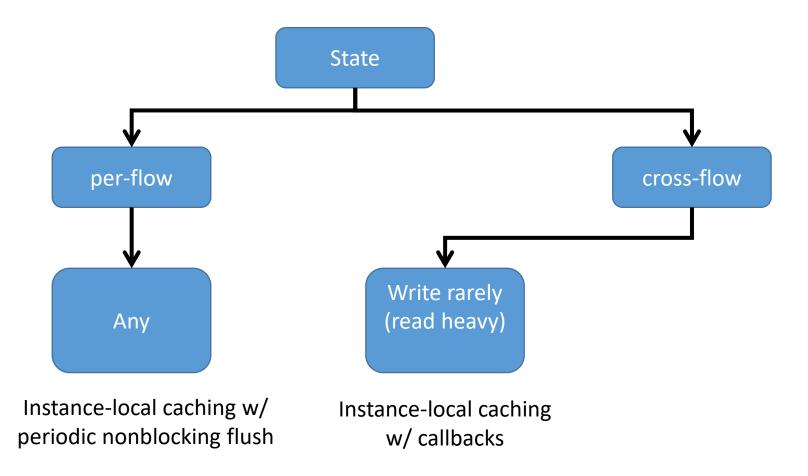
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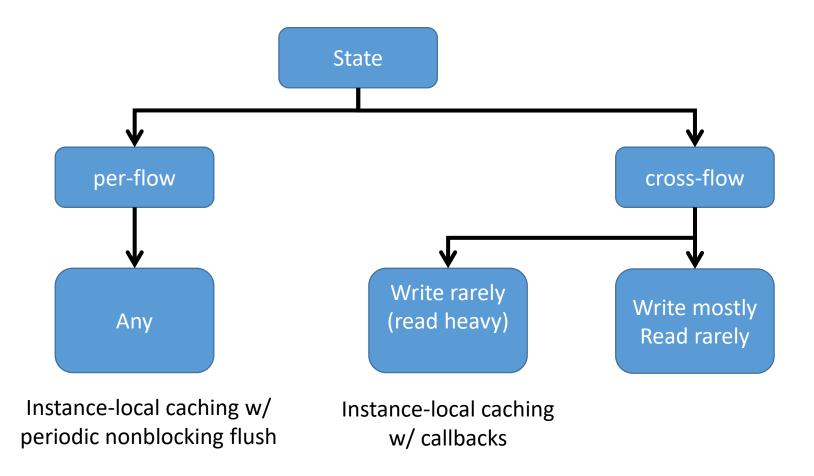
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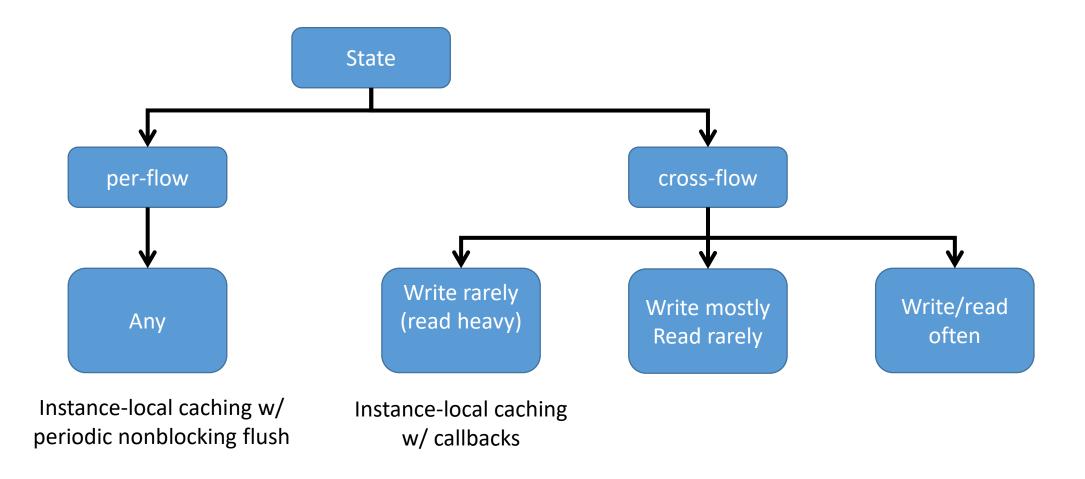


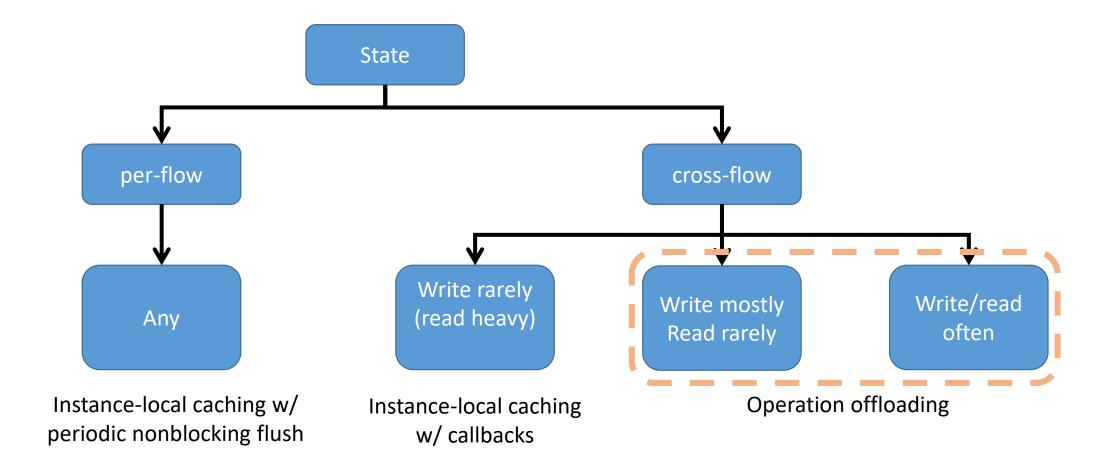


Instance-local caching w/ periodic nonblocking flush









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Operation	Description
Increment/Decrement a value	Increment or decrement the value stored at key by the given value
Push/pop a value to/from list	Push or pop the value in/from the list stored at the given key
Compare and update	Update the value, if the condition is true

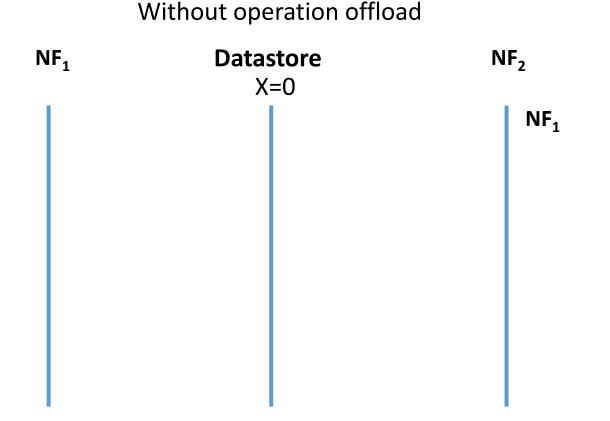
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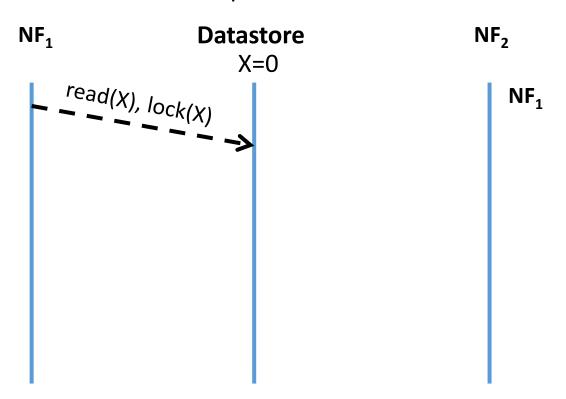
The datastore serializes operations issued by different instances for the same shared state object and applies them in the background

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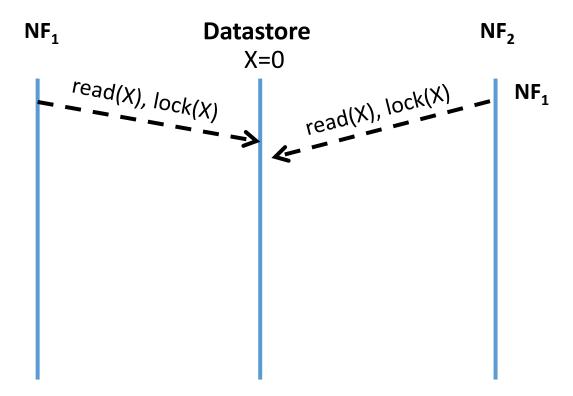
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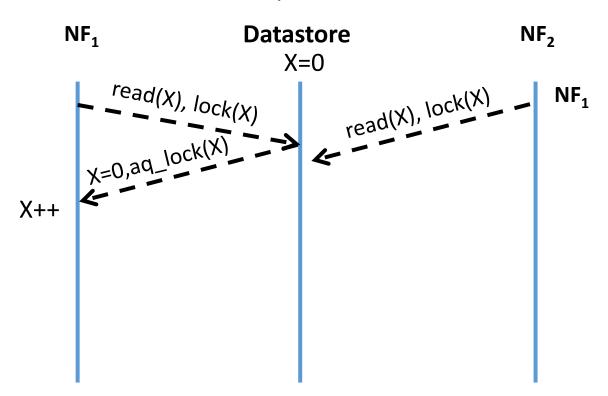
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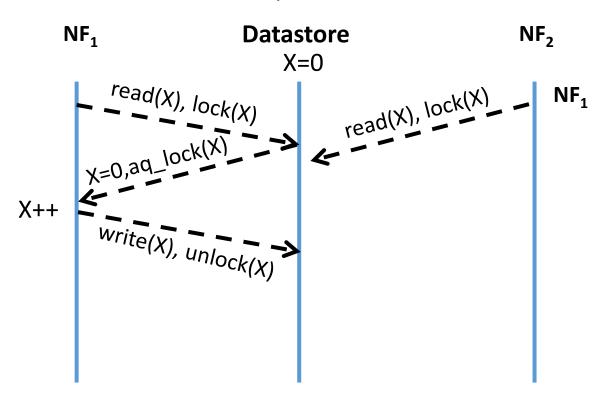
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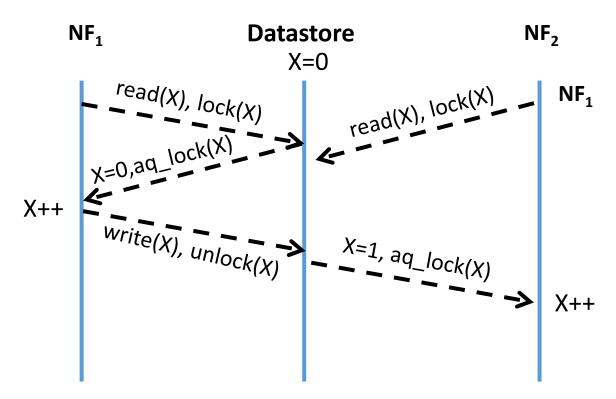
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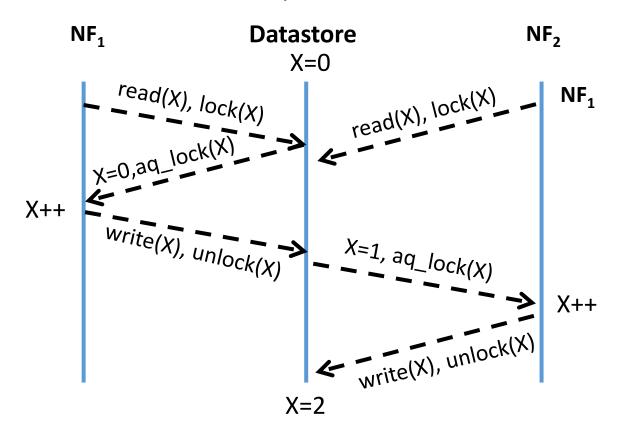
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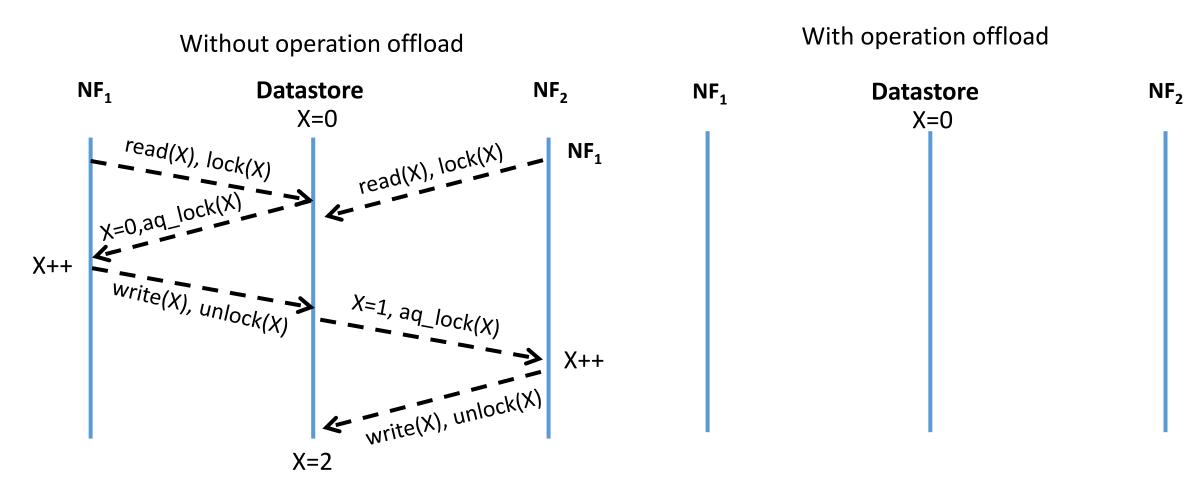
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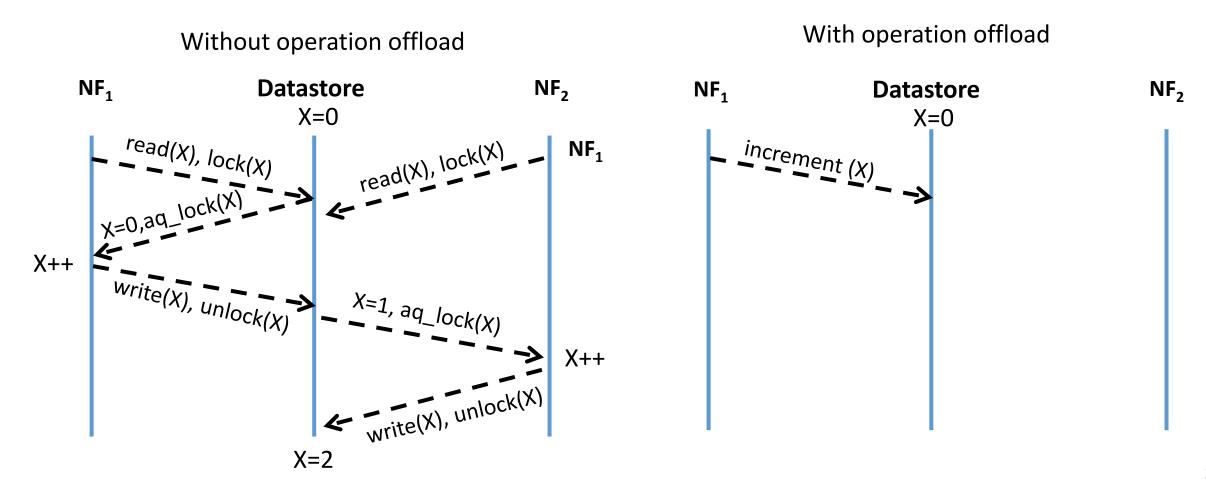


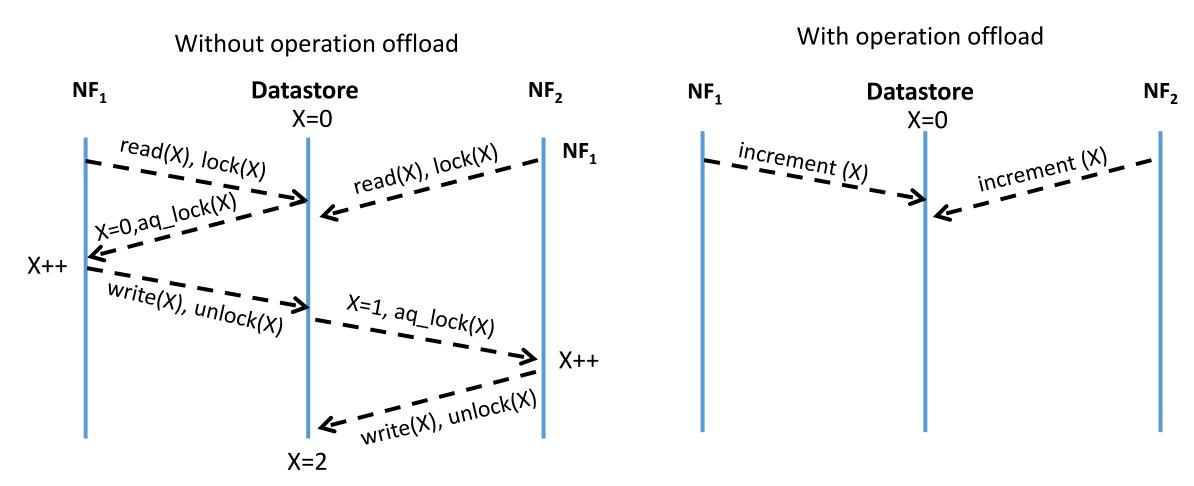
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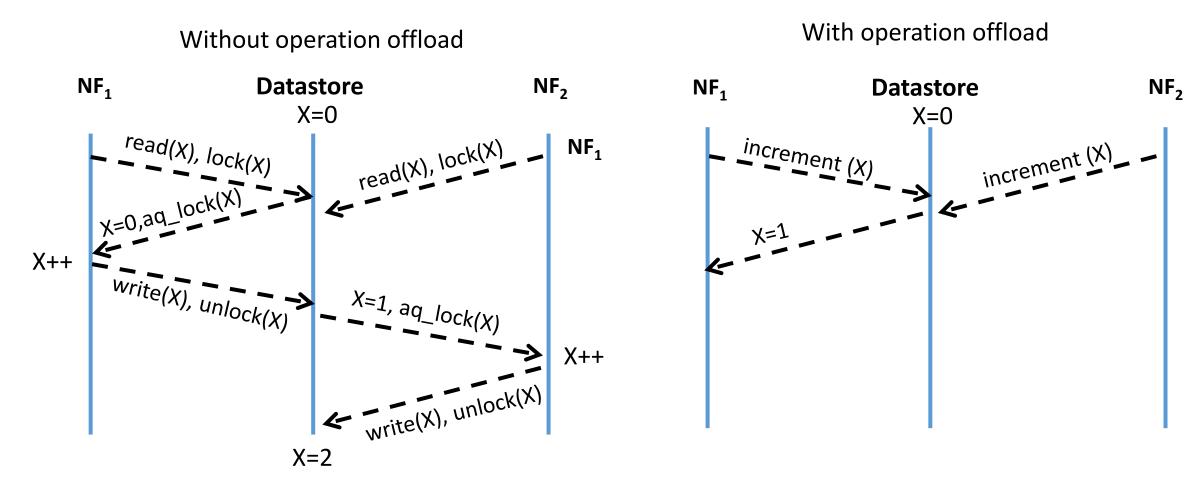


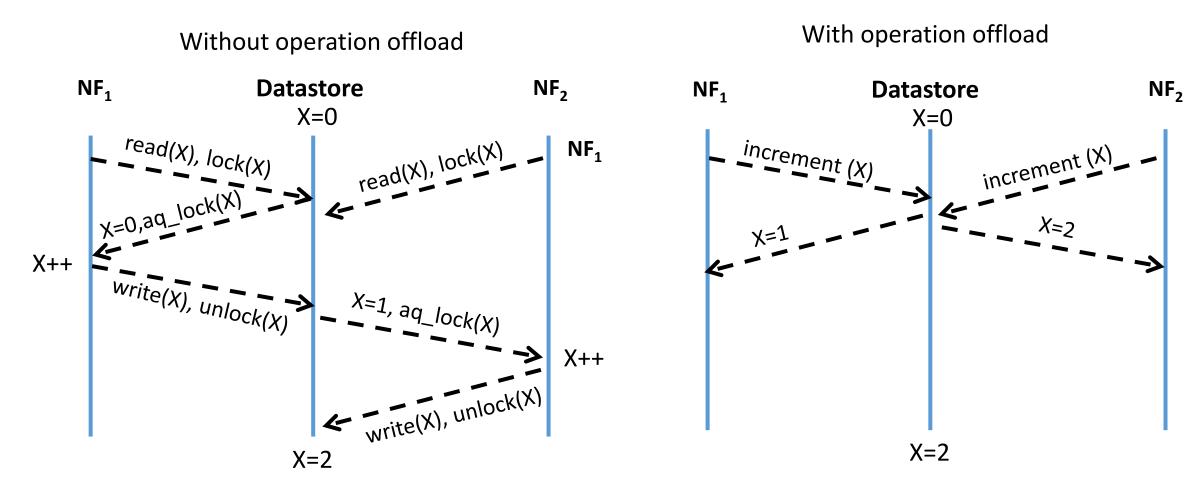
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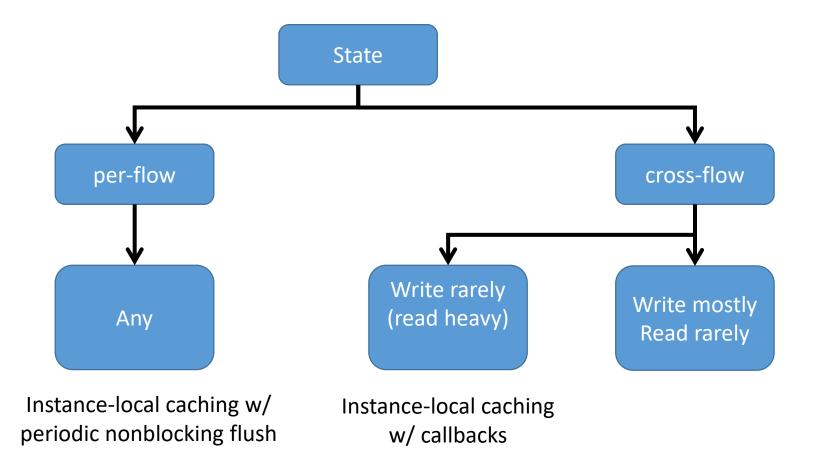




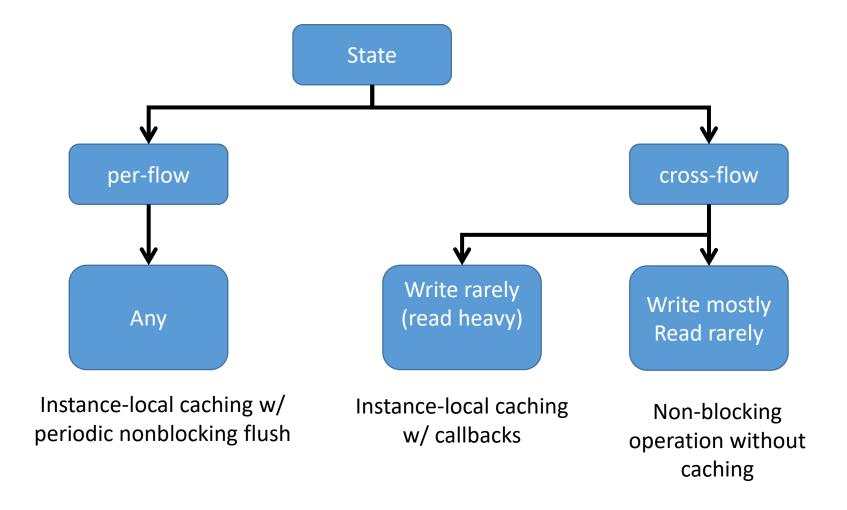




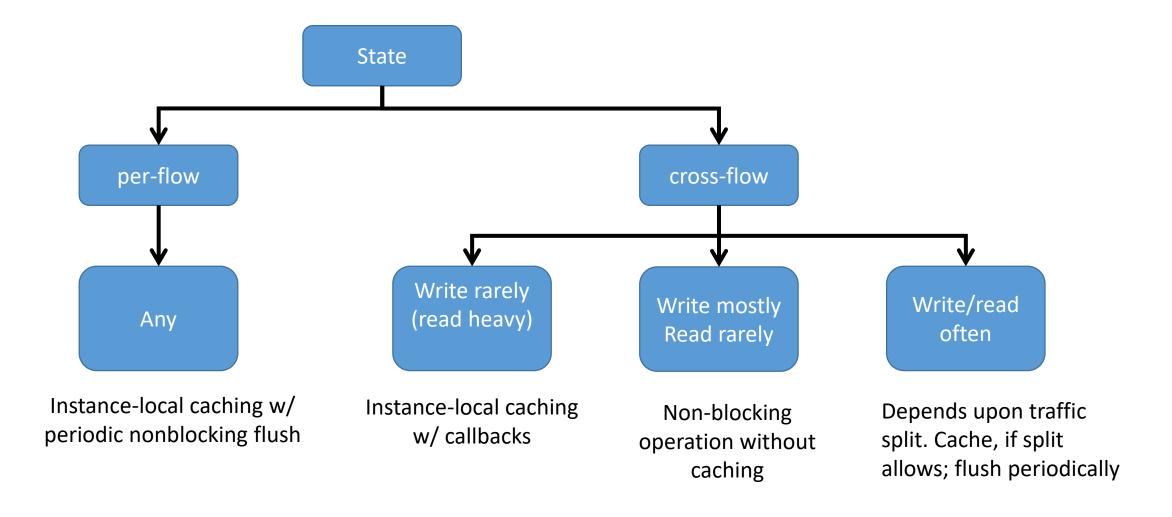
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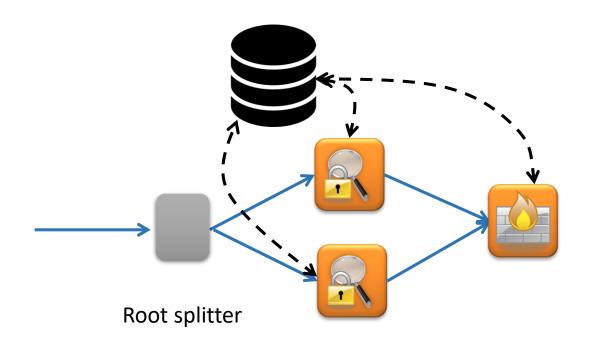
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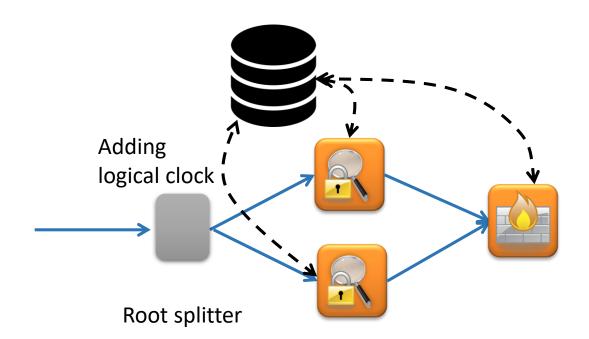
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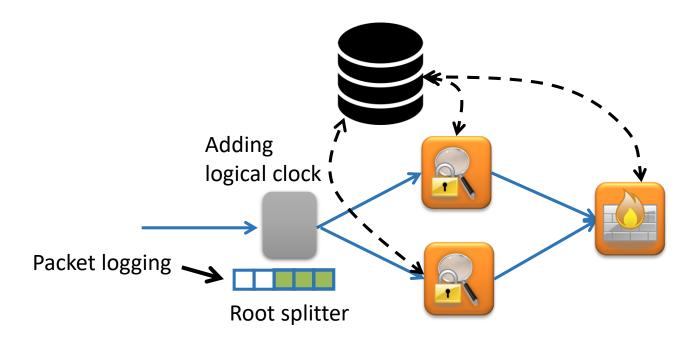
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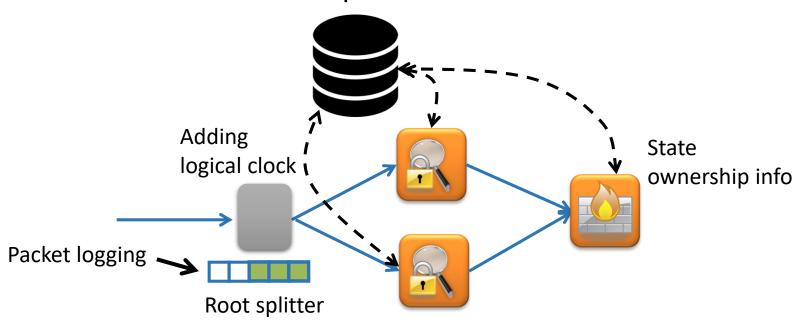
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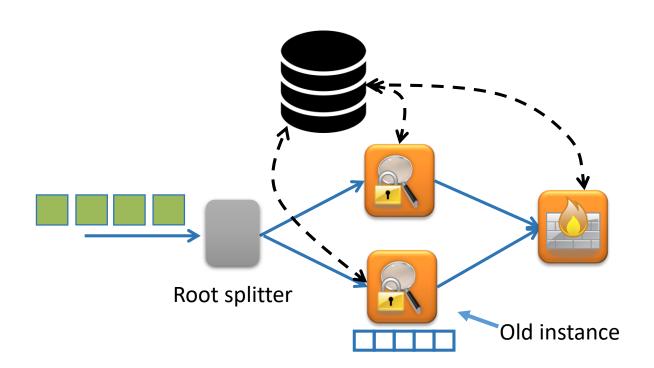


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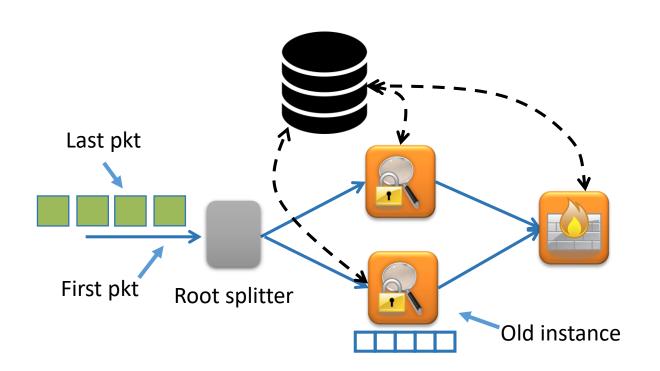
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CHC encodes state object's ownership information and logical clock associated with state operations as metadata

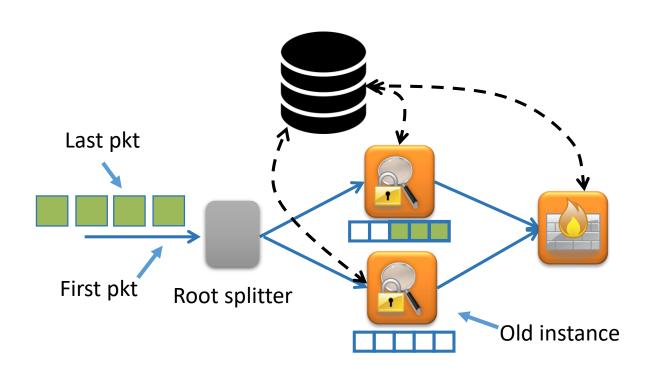




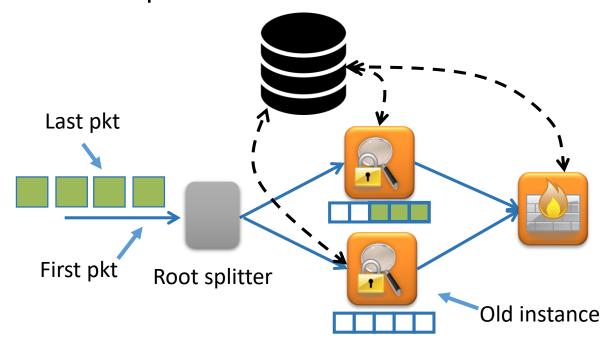
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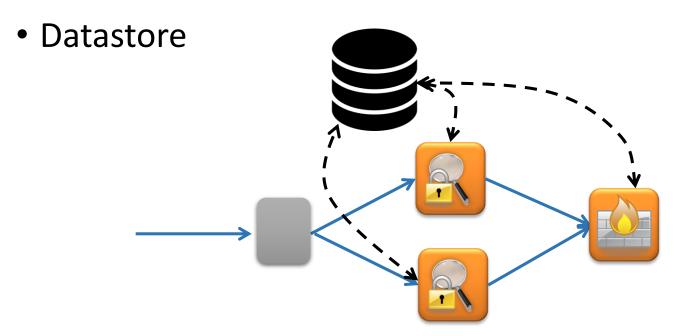


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- Cross-flow state does not require any special handling as operation offloading is used to update it



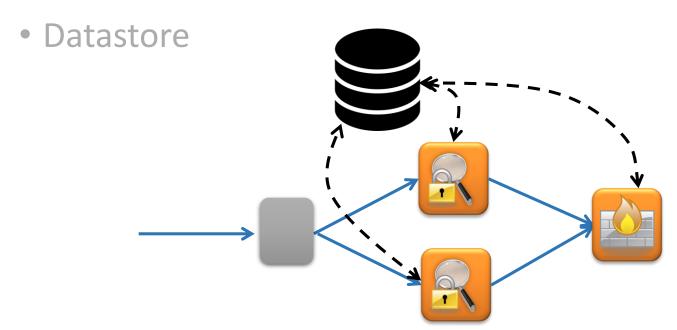
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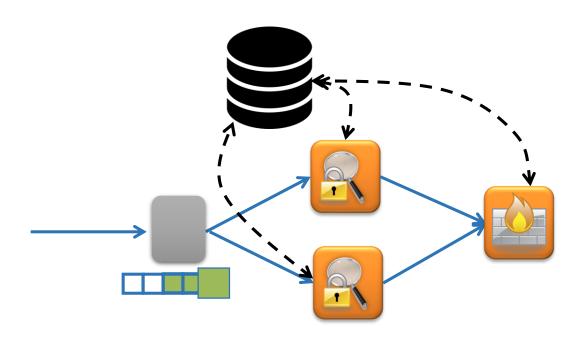
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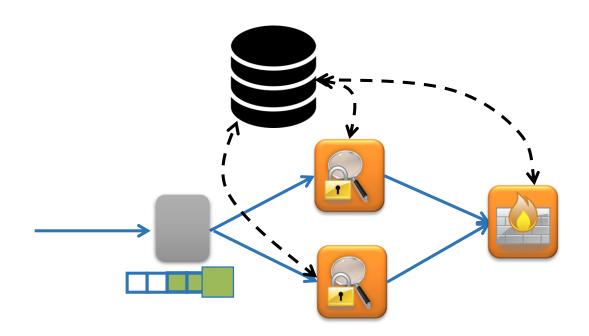
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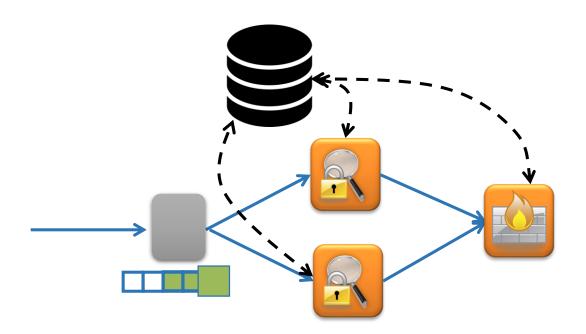


NF instance failure recovery:

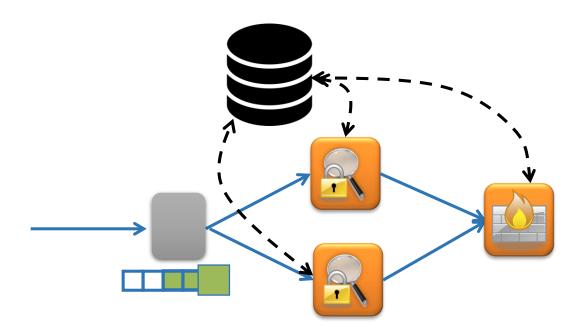
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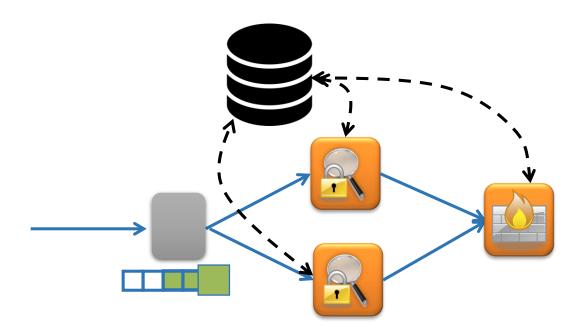
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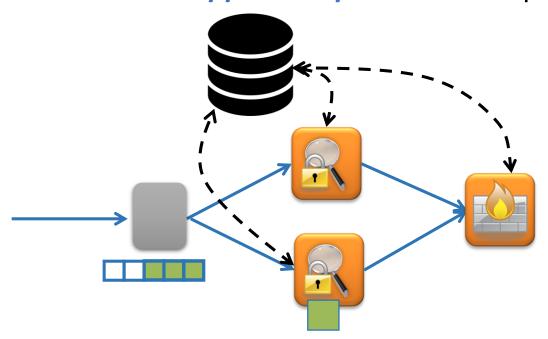
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- Root *replays* the packet

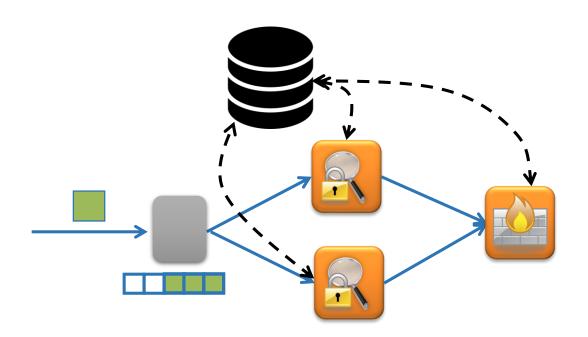


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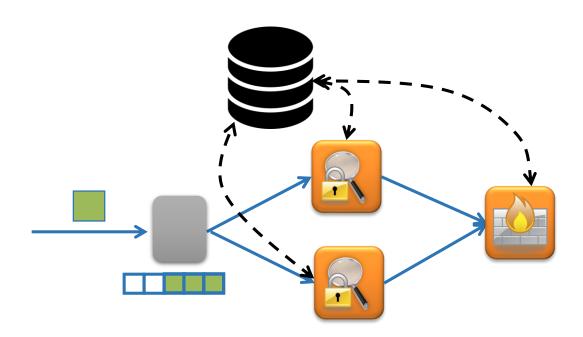


- Failover instance takes over
- Datastore associates the *failover instance ID* with the relevant state
- Root *replays* the packet
- Metadata is used to suppress duplicate state-update and processing

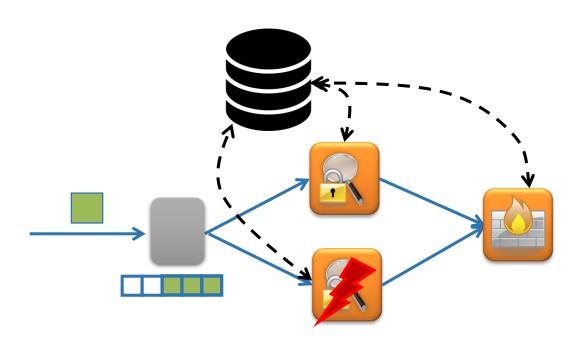




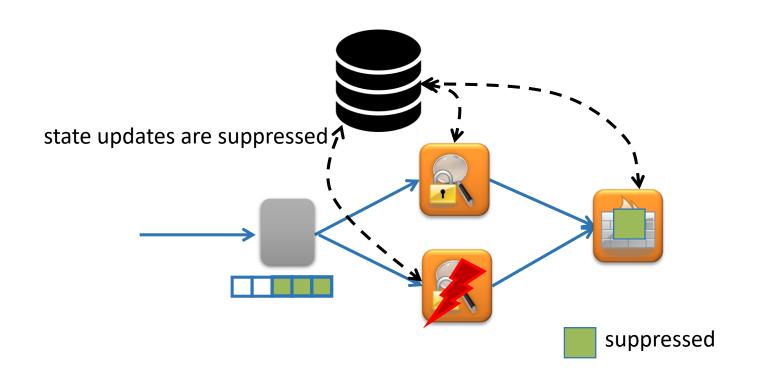
 Metadata (logical clocks) is used to suppress duplicate state updates at the datastore and duplicate packets at downstream NFs



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Implementation of CHC

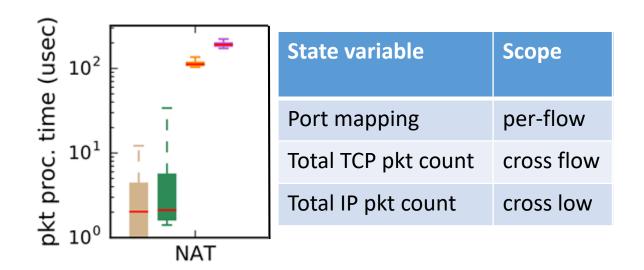
Implementation of CHC

- Prototype is implemented in C++
- Leverages Mellanox messaging accelerator for low latency communication

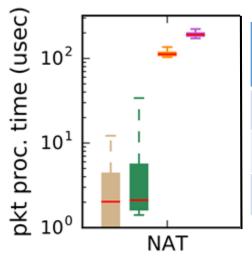
Implementation of CHC

- Prototype is implemented in C++
- Leverages Mellanox messaging accelerator for low latency communication

- We implemented four NFs on top of CHC
 - NAT
 - Trojan detector
 - Portscan detector
 - Load balancer

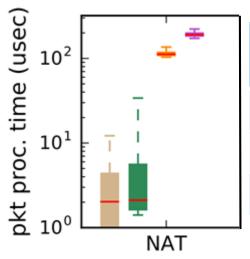


Traditional NF with infinite capacity



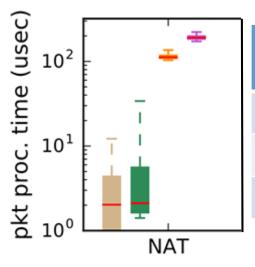
State variable	Scope	State Externalization
Port mapping	per-flow	
Total TCP pkt count	cross flow	
Total IP pkt count	cross low	

- Traditional NF with infinite capacity
- Externalized state operations



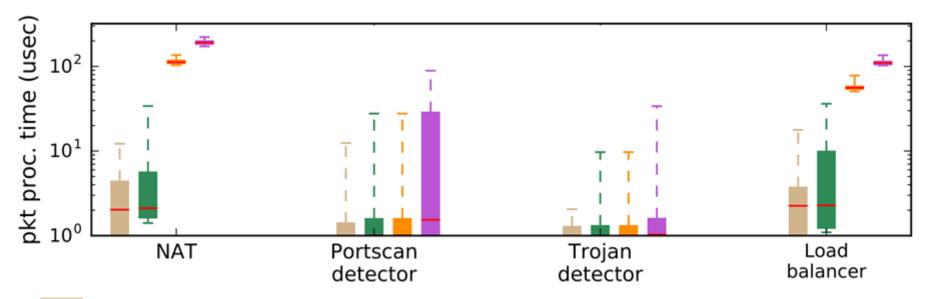
State variable	Scope	State Externalization	Caching
Port mapping	per-flow		
Total TCP pkt count	cross flow		
Total IP pkt count	cross low		

- Traditional NF with infinite capacity
- Externalized state operations
- State externalization with caching



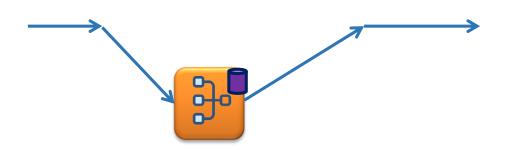
State variable	Scope	State Externalization	Caching	Asynch. + op offload
Port mapping	per-flow			
Total TCP pkt count	cross flow			
Total IP pkt count	cross low			

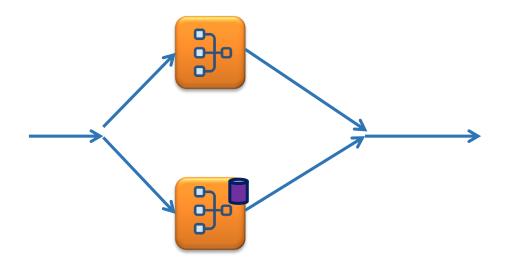
- Traditional NF with infinite capacity
- Externalized state operations
- State externalization with caching
- State externalization with caching and asynchronous + offloaded updates

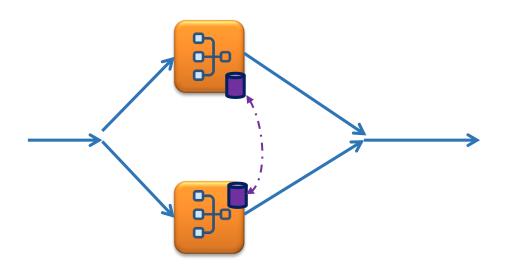


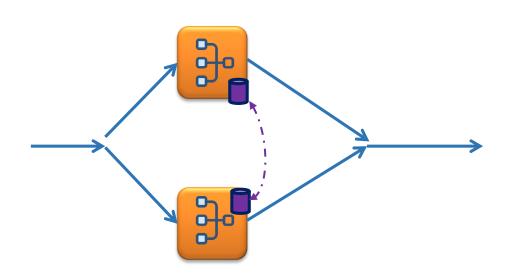
- Traditional NF with infinite capacity
- Externalized state operations
- State externalization with caching
- State externalization with caching and asynchronous + offloaded updates

Less than **0.6µs** increase in the median per-NF packet processing latency

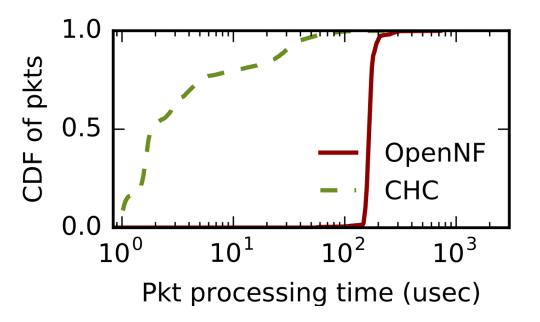


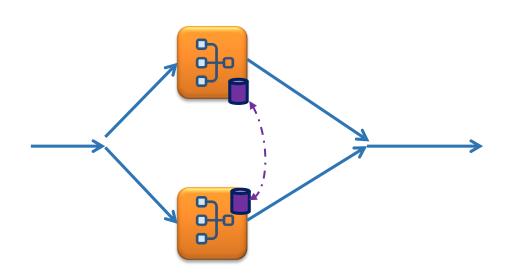




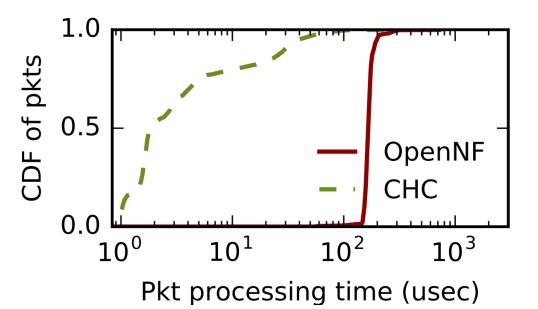


During cross instance state sharing

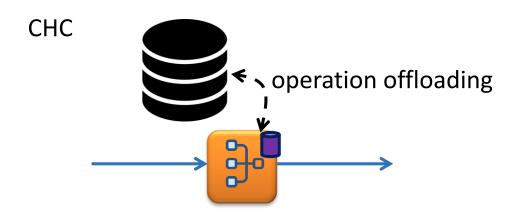


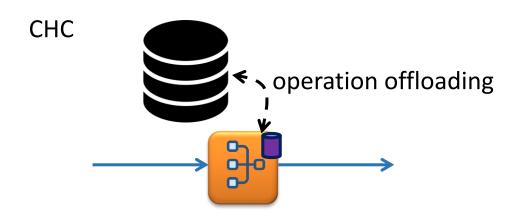


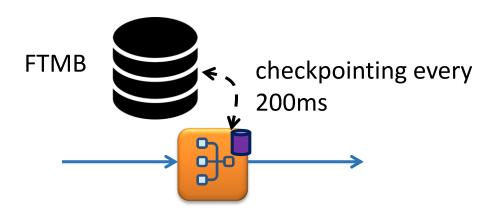
During cross instance state sharing

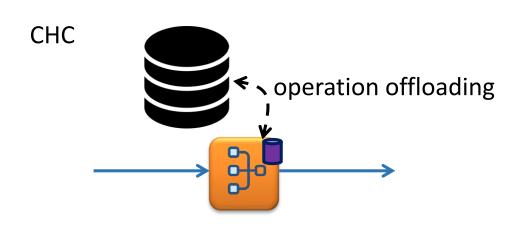


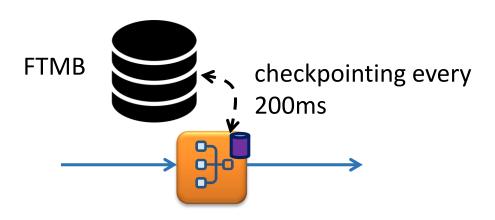
75th%-ile latency of CHC is **20 times lower than** OpenNF



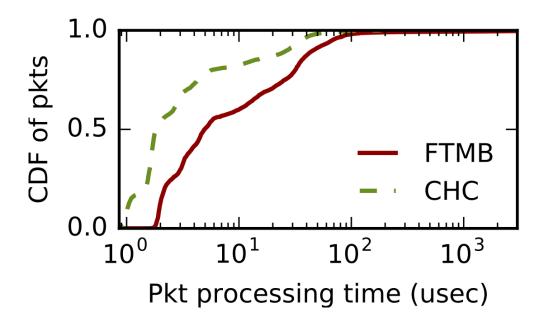


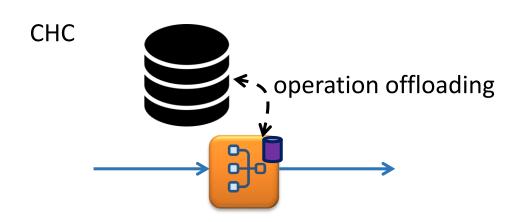


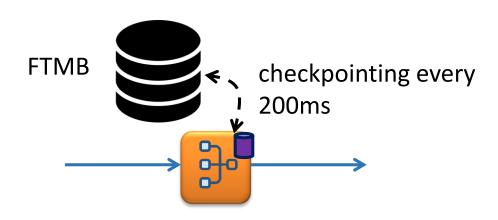




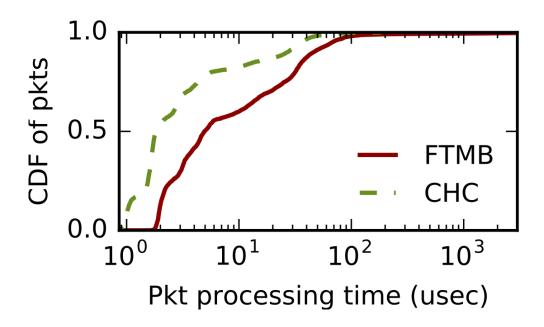
Ensuing Fault tolerance





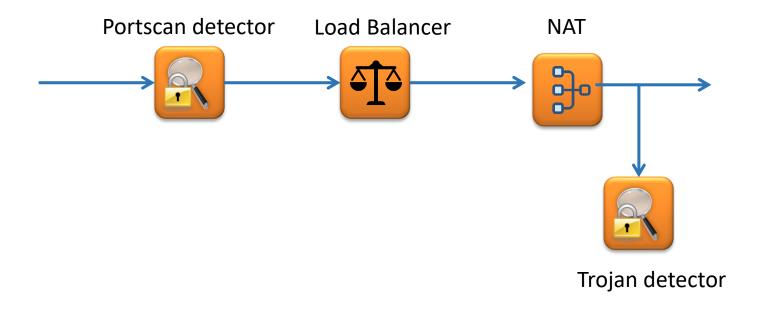


Ensuing Fault tolerance



75th%-ile latency of CHC is *6 times* lower than FTMB

Evaluation



CHC operates at line rate with an end-to-end median per packet processing overhead of *11.3us*

Evaluation

- State management performance
- Metadata overhead
- Correctness requirements:
 - State availability
 - Cross instance state transfer
 - Cross instance state sharing
 - Chain wide ordering
 - Duplication suppression
 - Fault tolerance

Summary

 CHC supports output equivalence and high performance state management for NFV chains

 It hides the complexity of handling states during dynamic actions (elastic scaling and failure recovery)

 It relies on managing state external to NFs, but couples it with several caching and state update algorithms to ensure low latency